fastppm: fast tumor phylogeny regression via tree-structured dual dynamic programming

Henri Schmidt*,1, Yuanyuan Qi*,2, Ben Raphael1, and Mohammed El-Kebir2



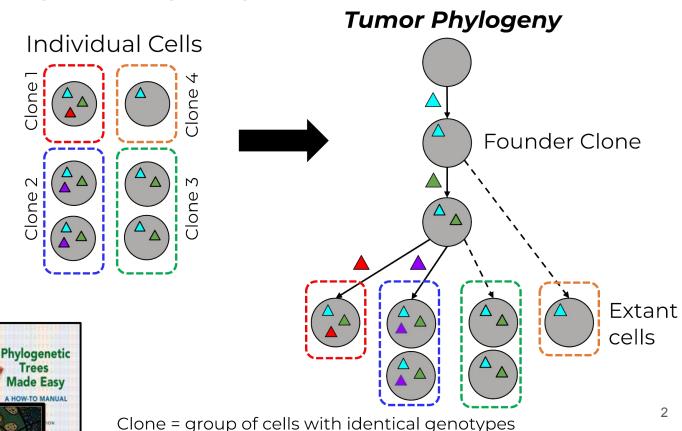


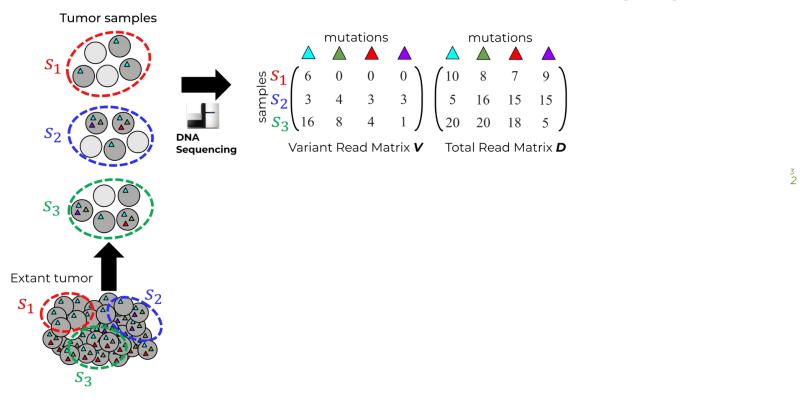
Reconstructing the evolutionary history of a tumor is a challenging and important open question

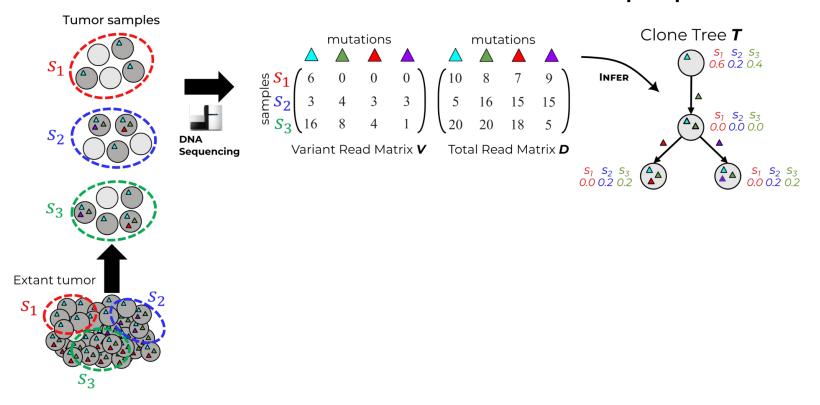
Molecular Evolution

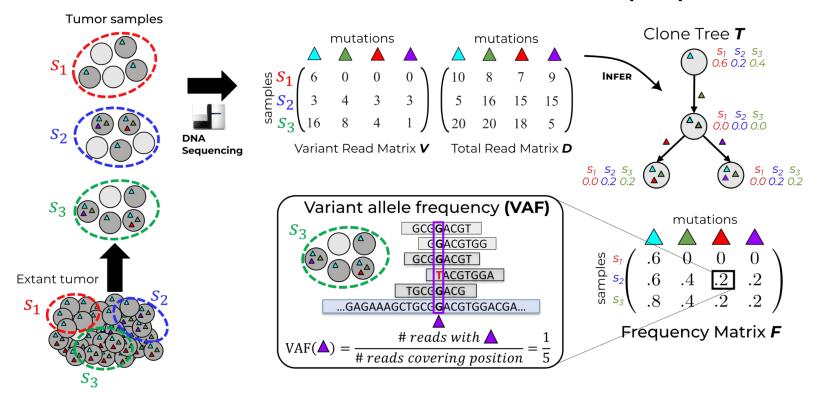
A Phylogenetic Approac

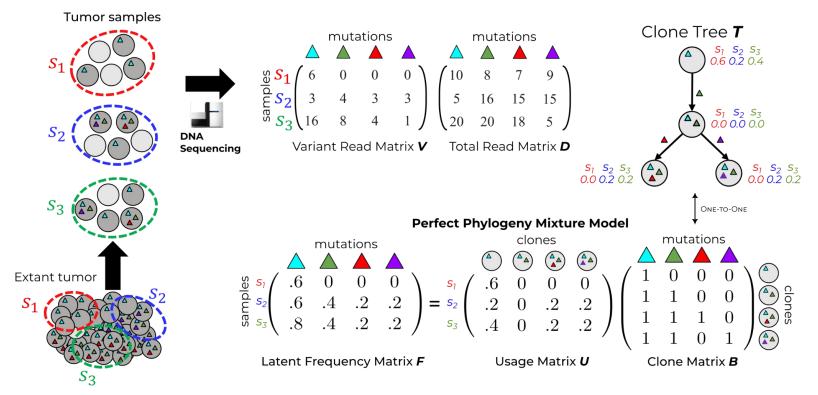
ENETICS



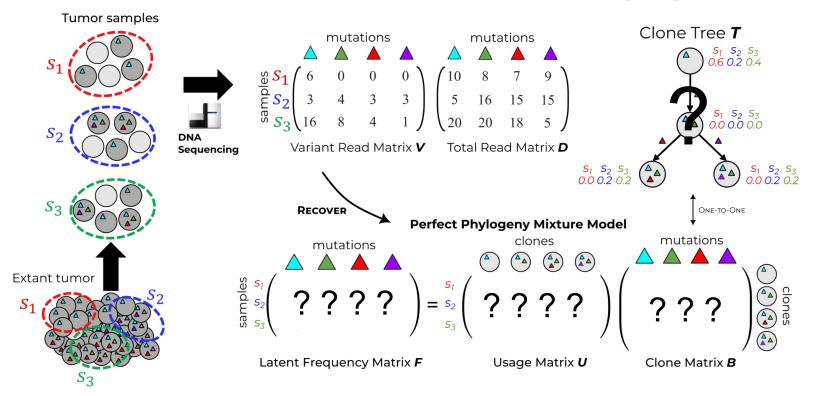




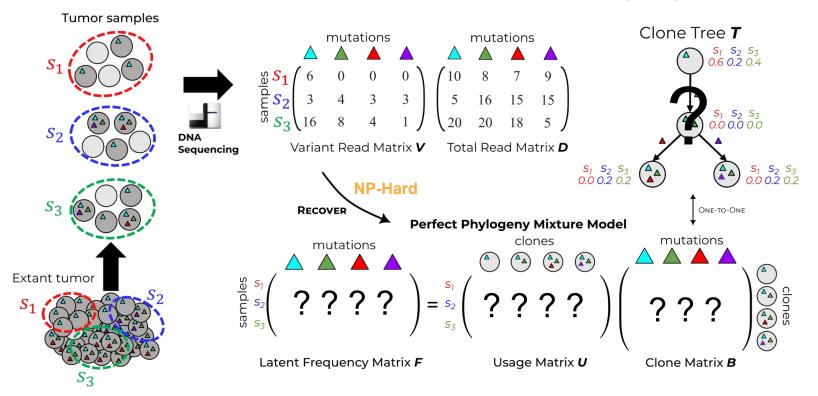




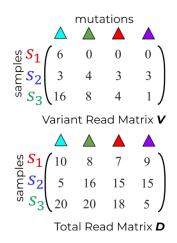
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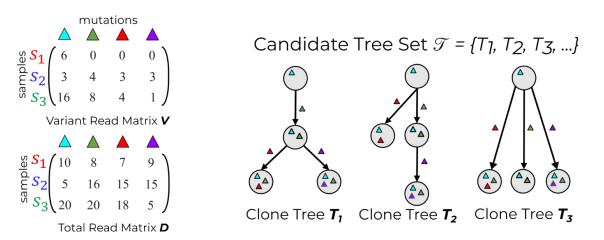


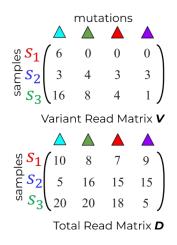
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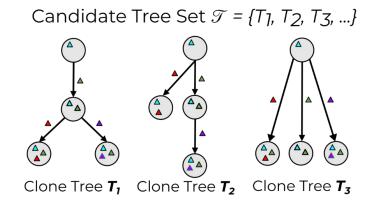


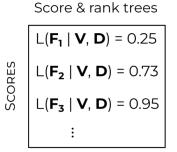
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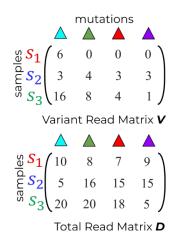


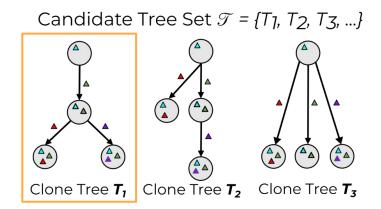






by fitting frequencies **F** to trees

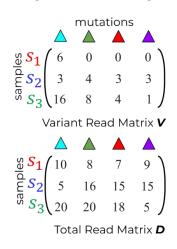


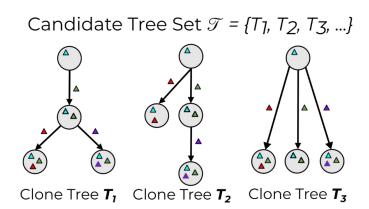


Select best tree(s)

Score & rank trees

by fitting frequencies **F** to trees





 $L(\mathbf{F_1} \mid \mathbf{V}, \mathbf{D}) = 0.25$ $L(\mathbf{F_2} \mid \mathbf{V}, \mathbf{D}) = 0.73$

Score & rank trees

 $L(\mathbf{F_2} \mid \mathbf{V}, \mathbf{D}) = 0.73$ $L(\mathbf{F_3} \mid \mathbf{V}, \mathbf{D}) = 0.95$ \vdots

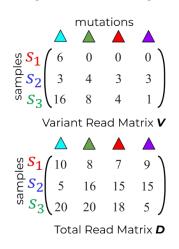
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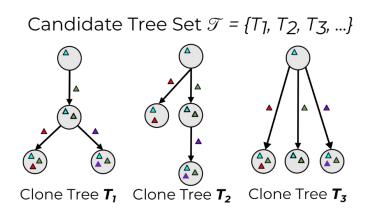
Trees are scored by **repeatedly*** solving the **perfect phylogeny regression problem**:

B is fixed

(PPR)
$$\min_{\mathbf{F},\mathbf{U}} \{ L(\mathbf{F} \mid \mathbf{V}, \mathbf{D}) : \mathbf{F} = \mathbf{U}\mathbf{B}, \mathbf{U} \ge 0, \mathbf{U}\mathbb{1} \le 1 \}$$

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Solving the PPR problem is the **key computational bottleneck** in phylogeny inference algorithms.

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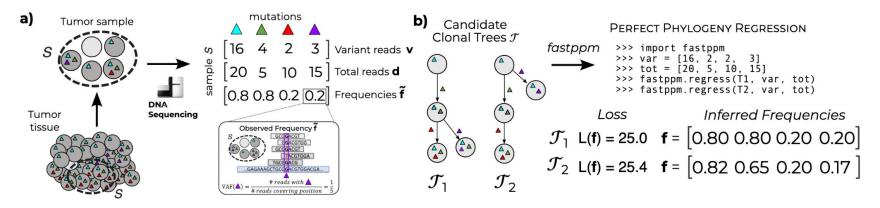
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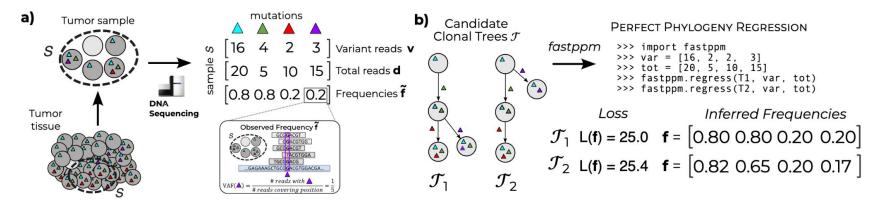
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Pitfall #2: General purpose solvers are slow

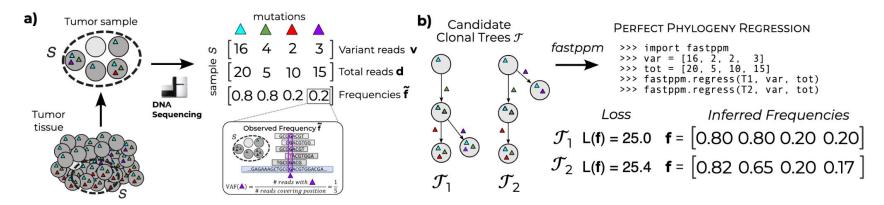


We introduce a new approach to the Perfect Phylogeny Regression problem, *fastppm*, using *tree structured dual dynamic programming (TSDDP)*.



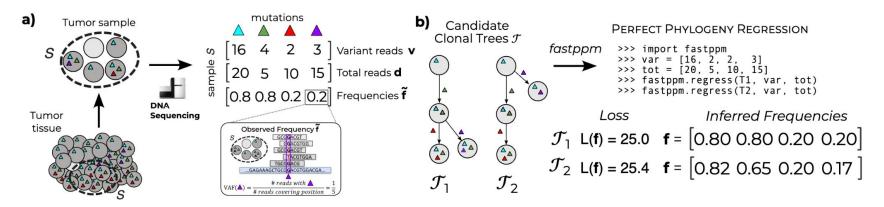
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- 2. *fastppm* is able to model arbitrary, **convex loss functions**, while maintaining its performance.

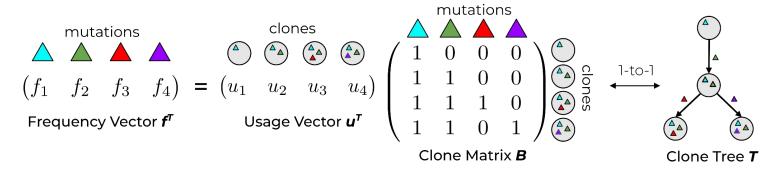


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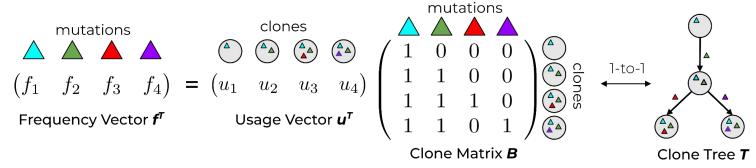
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On simulated data, replacing existing solvers with *fastppm* yields up to **400x** speed-ups and enables fast + accurate phylogenetic inference from **low-coverage** bulk DNA sequencing data.

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$$\min_{\mathbf{f},\mathbf{u}\in\mathbb{R}^n}\{\sum_{i=1}^n L_i(f_i): \mathbf{f}^T=\mathbf{u}^T\mathbf{B}, \mathbf{u}\geq 0, \mathbf{u}^T\mathbb{1}\leq 1\}$$

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mutations clones
$$(f_1 \quad f_2 \quad f_3 \quad f_4) = (u_1 \quad u_2 \quad u_3 \quad u_4)$$
Frequency Vector \mathbf{f}^T Usage Vector \mathbf{u}^T Clone Matrix \mathbf{B} Clone Tree \mathbf{T}

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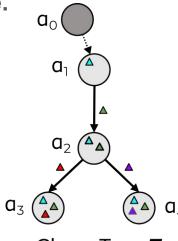
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In TSDDP, we first construct the **dual** optimization problem

(Dual-PPR)
$$\max_{\boldsymbol{\alpha} \in \mathbb{R}^{n+1}} \left\{ -\alpha_0 + \sum_{i=1}^n h_i (\alpha_i - \alpha_{\pi(i)}) : \boldsymbol{\alpha} \geq 0 \right\},$$

where $\pi(i)$ is the parent of vertex i in T and h_i is conjugate to L_i :





Clone Tree T

VARIABLES:

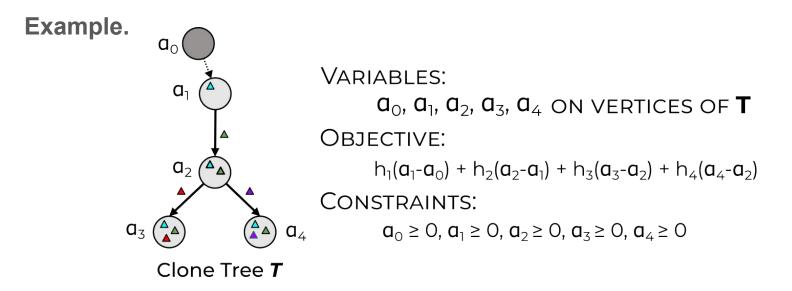
 a_0 , a_1 , a_2 , a_3 , a_4 on vertices of ${\bf T}$

OBJECTIVE:

$$h_1(a_1-a_0) + h_2(a_2-a_1) + h_3(a_3-a_2) + h_4(a_4-a_2)$$

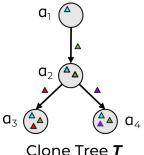
CONSTRAINTS:

$$a_0 \ge 0$$
, $a_1 \ge 0$, $a_2 \ge 0$, $a_3 \ge 0$, $a_4 \ge 0$



Key Idea: Solve the **dual** problem with a **bottom-up dynamic programming algorithm** over the clone tree **T**.

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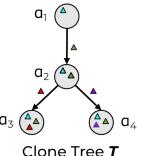
 $a_0 \ge 0$, $a_1 \ge 0$, $a_2 \ge 0$, $a_3 \ge 0$, $a_4 \ge 0$

Specifically, we define

$$J_i(\gamma) \triangleq \max_{\alpha > 0} \{ \sum_{j \in D(i)} h_j(\alpha_j - \alpha_{\pi(j)}) : \alpha_{\pi(i)} = \gamma \}$$

which is the optimal solution to the dual problem for the subtree rooted at vertex i, provided the dual variable of the parent of vertex i takes value y.

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which is the optimal solution to the dual problem for the subtree rooted at vertex i, provided the dual variable of the parent of vertex i takes value y. Then, this function satisfies the recurrence relation

$$J_i(\gamma) = \max_{\alpha_i > 0} \{ h_i(\alpha_i - \gamma) + \sum_{j \in \delta(i)} J_j(\alpha_i) \}$$
 (Recurrence Relation)

and TSDDP then computes the functions J_i in a **bottom-up** fashion.

For the *weighted* least squares loss, we solve the PPR problem in $\mathcal{Q}(n^{3/2}\log(\log(n)))$ time* over classes of random trees:

Efficient Projection onto the Perfect Phylogeny Model

Bei Jia* Surjyendu Ray jiabe@bc.edu Surjyendu Ray raysc@bc.edu Boston College

Sam Safavi safavisa@bc.edu jose.bento@bc.edu

Best known: $\mathcal{O}(n^2)$

Our result: $\mathcal{O}(n^{3/2}\log(\log(n)))$

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For the *piecewise linear loss* with k pieces, we solve the PPR problem in $\mathcal{O}(n\log^2(nk))$ time deterministically:

RESEARCH ARTICLE

A regression based approach to phylogenetic reconstruction from multi-sample bulk DNA sequencing of tumors

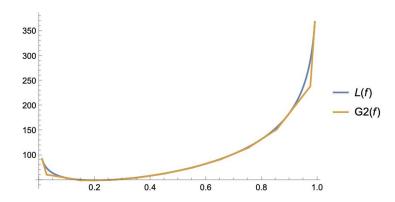
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Extensions to arbitrary, convex loss functions

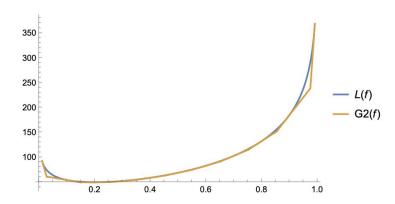
Approach #1: Piecewise Linear Approximation (*k*-PLA and PPLA)



Approximate one-dimensional convex loss function $L_i(f)$ with piecewise linear approximation using k pieces found via Taylor series expansion.

Extensions to arbitrary, convex loss functions

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Approximate one-dimensional convex loss function $L_i(f)$ with piecewise linear approximation using k pieces found via Taylor series expansion.

Approach #2: Structured Regression using Alternating Directions Method of Multipliers (ADMM)

ADMM:

```
\begin{array}{lll} x^{k+1} & := & \mathop{\rm argmin}_x L_\rho(x,z^k,y^k) & // \, x\text{-minimization} \\ z^{k+1} & := & \mathop{\rm argmin}_z L_\rho(x^{k+1},z,y^k) & // \, z\text{-minimization} \\ y^{k+1} & := & y^k + \rho(Ax^{k+1} + Bz^{k+1} - c) & // \, \, \text{dual update} \end{array}
```

Using ADMM, we reduce solving the Perfect Phylogeny Regression problem for **arbitrary** convex loss functions to a sequence of L_2 subproblems.

Fast regression under the perfect phylogeny model

Results (L₂ / least squares loss):

- fastppm achieved a 40-125x speed up over the next best performing method projectppm.
- All methods achieved the exact same objective value on all instances.
- Blackbox convex optimization solvers were significantly slower than projectppm and fastppm.
- Excludes model build time which is in practice non-negligible for blackbox solvers.

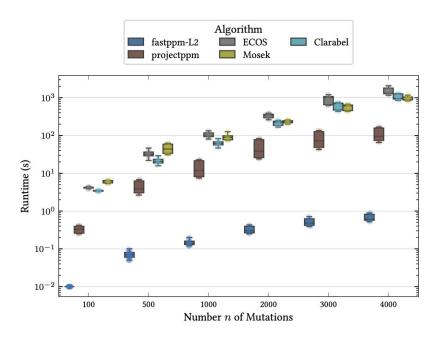


Figure: Runtime of existing methods for the Perfect Phylogeny Regression problem when varying the number of clones/mutations.

Improving existing phylogeny inference methods with fastppm

We replaced calls to existing perfect phylogeny regression algorithms in Sapling, CITUP, and Orchard with calls to *fastppm*:

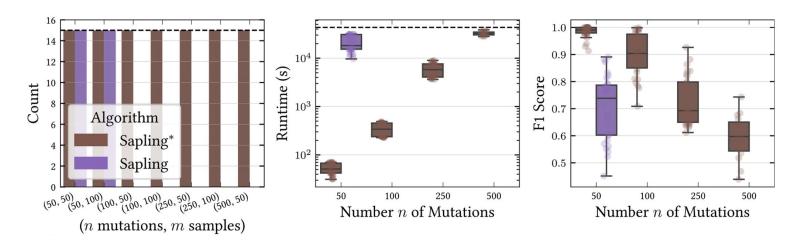


Figure: Number of successful instances within a twelve hour time limit, runtime, and accuracy of Sapling compared to Sapling* on simulated data.

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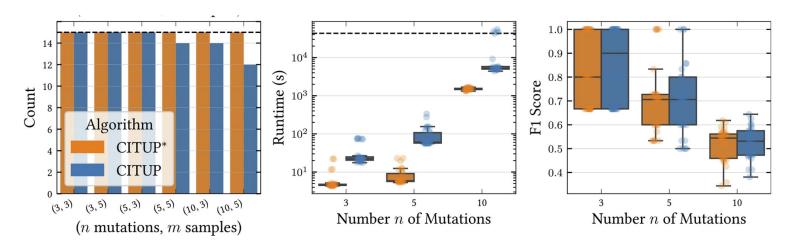


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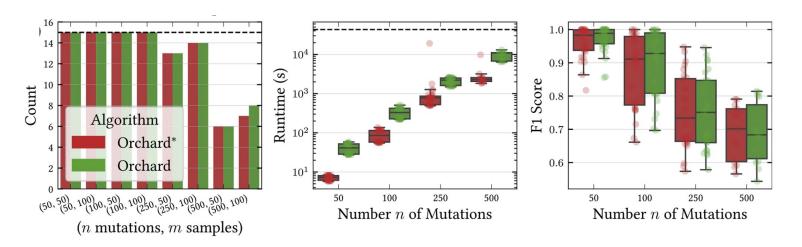
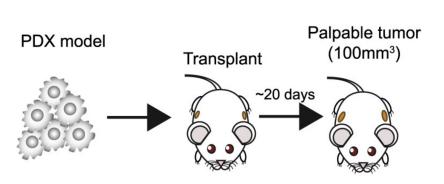


Figure: Number of successful instances within a twelve hour time limit, runtime, and accuracy of Orchard compared to Orchard* on simulated data.

fastppm improves frequency estimation in low coverage settings

Downsampled reads from a patient derived xenograft (POP66) is a mouse model of colorectal cancer (n = 64 mutations, m = 8 samples, 50x coverage):



Data from: (Rehman et al. 2021)

Method	Metric	Objective
Orchard* Orchard Sapling*	$egin{aligned} \ ilde{\mathbf{F}} - \hat{\mathbf{F}}\ _F^2 \ \ ilde{\mathbf{F}} - \hat{\mathbf{F}}\ _F^2 \ \ ilde{\mathbf{F}} - \hat{\mathbf{F}}\ _F^2 \end{aligned}$	3.092 2.448 2.181
Orchard* Orchard Sapling*	$egin{aligned} -\log \mathbb{P}(\mathbf{V} \mid \mathcal{T}, \hat{\mathbf{F}}, \mathbf{D}) \ -\log \mathbb{P}(\mathbf{V} \mid \mathcal{T}, \hat{\mathbf{F}}, \mathbf{D}) \ -\log \mathbb{P}(\mathbf{V} \mid \mathcal{T}, \hat{\mathbf{F}}, \mathbf{D}) \end{aligned}$	10790.9 10793.5 10720.6

We applied Orchard, Orchard*, which use the L₂ loss, to Sapling*, which uses the binomial loss, to recover the clonal tree and mutation frequencies.

Thank You

Collaborators

Yuanyuan Qi (Co-first Author) Mohammed El-Kebir Ben Raphael

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Hirak Sarkar Henri Schmidt

Richard Zhang Ahmed Shuaibi

Peter Halmos Akhil Jakatdar

Yihang Shen Gary Hu

W. Howard-Synder Clover Zheng

Michael Wilson Viola Chen

Julian Gold



The Raphael Lab



fastppm is implemented in C++ and is available on GitHub



Mohammed El-Kebir



Manuscript is accessible through *Bioinformatics*

SCRATCH

However, existing methods for the PPR problem are flawed

1. **Do not directly** model the read count data, which e.g. hinders analysis of **low-coverage DNA sequencing**

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Example 1.

State-of-the-art phylogeny inference pipelines *CITUP*, *AncesTree*, *CALDER*, *Pairtree*, *Orchard*, and *fastBE* use the following two loss functions:

$$L_1(\mathbf{F}, \mathbf{V}, \mathbf{D}) = \sum_{i=1}^m \sum_{j=1}^n |f_{ij} - \tilde{f}_{ij}| \text{ where } \tilde{f}_{ij} = v_{ij}/d_{ij}$$

$$L_2(\mathbf{F}, \mathbf{V}, \mathbf{D}) = \sum_{i=1}^m \sum_{j=1}^n w_{ij} (f_{ij} - \tilde{f}_{ij})^2 \text{ where } \tilde{f}_{ij} = v_{ij}/d_{ij}, w_{ij} \ge 0$$
Observed frequency (VAF)

which do not directly model the read count data, instead collapsing it to a frequency.

However, existing methods for the PPR problem are flawed

- Do not directly model the read count data, which e.g. hinders analysis of low-coverage DNA sequencing
- 2. Employ *slow*, black-box convex optimization software which *do not exploit the structure of the regression problem*

Example 2.

In contrast, phylogeny inference pipelines (e.g. Sapling, PhyloWGS) which model observations using the probabilistic read-count model v_{ij} Binomial(f_{ij} , d_{ij}), e.g.,

$$L_{\mathrm{Bin}}(\mathbf{F},\mathbf{V},\mathbf{D}) = -\sum_{i=1}^{m} \sum_{j=1}^{n} [v_{ij} \log f_{ij} + (d_{ij} - v_{ij}) \log (1 - f_{ij})].$$

must resort to blackbox convex optimization software which is prohibitively slow.

Then, we solve the **dual** problem with a **bottom-up dynamic programming algorithm over T**.

- (i) Fix a representation $\mathcal{R}(J_i)$ for each J_i .
- (ii) Compute the representation $\mathcal{R}(J_i)$ at the leaf nodes.
- (iii) Compute the representation $\mathcal{R}(J_i)$ at a node i provided the representations $\mathcal{R}(J_j)$ at all children $j \in \delta(i)$.
- (iv) Solve the one-dimensional optimization problem $\max_{\alpha_0 \geq 0} \{-\alpha_0 + J_r(\alpha_0)\}$ using the representation of the root node $\mathcal{R}(J_r)$.

The TSDDP Algorithm